Mathematical Principles in Visual Computing: Solving Polynomial Systems
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Many Computer Vision problems can be solved by finding the roots of a polynomial system:

- camera pose estimation from point correspondences;
- camera relative motion estimation from point correspondences;
- image distortion calibration;
- point triangulation;
- ...

Polynomial Systems in Computer Vision
Solving Polynomial Systems

- no general method;
- several mathematical tools exist. For a given problem, a tool can be more adapted than the others.
introduced in 1965 by Bruno Buchberger (now at the Johannes Kepler University in Linz) in his Ph.D. thesis (named after his advisor Wolfgang Gröbner) to study sets of polynomials
A Polynomial System

Let consider the following polynomial system:

\[
\begin{align*}
L_1 & \quad 2x^2 + y^2 - 2z + 3z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2y^2 + y^2z^2 - 2 = 0
\end{align*}
\]

Hint: try to remove \( x \) from the first equation

Replace \( L_1 \) by \( L_1 - 2L_2 \):

\[
\begin{align*}
L'_1 & \quad y^2 - 4z + z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2y^2 + y^2z^2 - 2 = 0
\end{align*}
\]
A Real Polynomial System (continued)

\[
\begin{align*}
L_1' & \left\{ y^2 - 4z + z^2 + 5 = 0 \right. \\
L_2 & \left\{ x^2 + z + z^2 = 0 \right. \\
L_3 & \left\{ x^2y^2 + y^2z^2 - 2 = 0 \right. \\
L_4 & \left\{ y^2z + 2 = 0 \right. \\
\end{align*}
\]

Hint: try to remove \(x\) from the second equation:

Adding \(y^2L_2 - L_3\):

\[
\begin{align*}
L_1' & \left\{ y^2 - 4z + z^2 + 5 = 0 \right. \\
L_2 & \left\{ x^2 + z + z^2 = 0 \right. \\
L_3 & \left\{ x^2y^2 + y^2z^2 - 2 = 0 \right. \\
L_4 & \left\{ y^2z + 2 = 0 \right. \\
\end{align*}
\]
A Real Polynomial System (continued)

\[
\begin{align*}
L'_1 & \quad y^2 - 4z + z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2 y^2 + y^2 z^2 - 2 = 0 \\
L_4 & \quad y^2 z + 2 = 0
\end{align*}
\]

Hint: try to remove \( y \) from the first equation

Add \( zL'_1 - L_4 \):

\[
\begin{align*}
L'_1 & \quad y^2 - 4z + z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2 y^2 + y^2 z^2 - 2 = 0 \\
L_4 & \quad y^2 z + 2 = 0 \\
L_5 & \quad 5z - 4z^2 + z^3 - 2 = 0
\end{align*}
\]
A Real Polynomial System (continued)

\[
\begin{aligned}
L_1' & \quad y^2 - 4z + z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2 y^2 + y^2 z^2 - 2 = 0 \\
L_4 & \quad y^2 z + 2 = 0 \\
L_5 & \quad 5z - 4z^2 + z^3 - 2 = 0
\end{aligned}
\]

Hint: \( L_5 \) is a polynomial in \( z \) only

\[5z - 4z^2 + z^3 - 2 = (z - 1)^2(z - 2)\]

Each possible value for \( z \) gives a new polynomial system in \( x \) and \( y \) only.
Solving a Univariate Polynomial

- closed form up to degree 4;
- for higher degrees:
  - the companion matrix method: The *companion matrix* of
    \[ p(z) = z^n + a_{n-1}z^{n-1} + \ldots + a_1z + a_0 \]
    is
    \[
    C = \begin{bmatrix}
    0 & \cdots & 0 & -a_0 \\
    1 & 0 & \cdots & -a_1 \\
    & 1 & 0 & \cdots & -a_2 \\
    & & \ddots & \vdots & \ddots \\
    & & & 1 & -a_{n-1} \\
    \end{bmatrix}.
    \]
    Its eigenvalues are the roots of \( p(z) \) (because \( p(z) \) is the
    characteristic polynomial \( \det(zI - C) \) of \( C \)).
  - Sturm’s bracketing method (slightly less stable but much faster).
Two Gröbner bases

\[
\begin{align*}
L'_1 & \quad y^2 - 4z + z^2 + 5 = 0 \\
L_2 & \quad x^2 + z + z^2 = 0 \\
L_3 & \quad x^2 y^2 + y^2 z^2 - 2 = 0 \\
L_4 & \quad y^2 z + 2 = 0 \\
L_5 & \quad 5z - 4z^2 + z^3 - 2 = 0
\end{align*}
\]

\[
\{ y^2 - 4z + z^2 + 5, x^2 + z + z^2, x^2 y^2 + y^2 z^2 - 2, y^2 z + 2, 5z - 4z^2 + z^3 - 2 \}
\]
is a Gröbner basis.

\[
\{ y^2 - 4z + z^2 + 5, x^2 + z + z^2, 5z - 4z^2 + z^3 - 2 \}
\]
is also a Gröbner basis.
A Gröbner basis is a set of polynomials \( \{g_1, \ldots, g_t\} \), such that the system

\[
\begin{align*}
g_1(x_1,\ldots,x_n) &= 0 \\
#\ldots\
g_t(x_1,\ldots,x_n) &= 0
\end{align*}
\]

has the same solutions as the original one, but with some specific properties that make the new system easier to solve than the original one, OR AT LEAST USEFUL to solve the original one.
We can create new equations from:

▶ linear combinations of existing equations. Gauss-Jordan elimination algorithm to simplify the system.

\[
\begin{align*}
2x^2 + xy + y^2 + 1 &= 0 \\
x^2 - xy + 2y^2 - 1 &= 0
\end{align*}
\]

in matrix form:

\[
\begin{bmatrix}
2 & 1 & 1 & 1 \\
1 & -1 & 2 & -1
\end{bmatrix}
\begin{bmatrix}
x^2 \\
x y \\
y^2 \\
1
\end{bmatrix} = 0.
\]

\[
\begin{bmatrix}
1 & 0 & 1 & 0 \\
0 & 1 & -1 & 1
\end{bmatrix}
\begin{bmatrix}
x^2 \\
x y \\
y^2 \\
1
\end{bmatrix} = 0.
\]

▶ *algebraic* combinations of existing equations.
We can create new equations from:

- linear combinations of existing equations.
- *algebraic* combinations of existing equations.
- the remainder of polynomial divisions (used by Buchberger’s algorithm).
Monomials

Definition. A monomial in \(x_1, \ldots, x_n\) is a product of the form:

\[x_1^{\alpha_1} \cdot x_2^{\alpha_2} \cdots x_n^{\alpha_n},\]

where all the exponents \(\alpha_1, \ldots, \alpha_n\) are nonnegative integers, sometimes noted \(x^\alpha\) with \(\alpha = (\alpha_1, \ldots, \alpha_n)\).

Examples: \(x, x^2, x^2y, x^2yz^3\)
Definition. A polynomial $f$ in $x_1, \ldots, x_n$ with coefficients in a field $k$ is a finite linear combination with coefficients in $k$ of monomials. A polynomial is written in the form

$$f = \sum_\alpha a_\alpha x^\alpha, \quad a_\alpha \in k$$

with

- $a_\alpha$ the coefficient of the monomial $x^\alpha$.
- If $a_\alpha \neq 0$, then we call $a_\alpha x^\alpha$ a term of $f$. 
A monomial order is any relation on the set of monomials $x^\alpha$ in $k[x_1, \ldots, x_n]$ satisfying:

1. $>$ is a total (linear) ordering relation:
   there is only one possible to order in increasing order under $>$
   a set of monomials;

2. $>$ is compatible with multiplication:
   if $x^\alpha > x^\beta$ and $x^\gamma$ is any monomial, then
   $x^\alpha x^\gamma = x^{\alpha+\gamma} > x^\beta x^\gamma = x^{\beta+\gamma}$;

3. $>$ is a well-ordering:
   every nonempty set of monomials has a smallest element
   under $>$. 

Monomial Order
Monomial Order on $k[x]$

The only monomial order on $k[x]$ is the degree order, given by:

$$\ldots > x^{n+1} > x^n > \ldots > x^2 > x > 1.$$
Monomial Orders on $k[x_1, \ldots, x_n]$

For polynomials in several variables, there are many choices of monomial orders.

Let’s first define an order on the variables: $x_1 > x_2 > \ldots > x_n$ (this is not a monomial order), and $x > y > z$. 
Monomial Orders on $k[x_1, \ldots, x_n]$ - the Lexicographic Order $>_{\text{lex}}$

**Definition.** The lexicographic order: analogous to the ordering of words in a dictionary.

For example, under this order $>_{\text{lex}}$:

$$x^2 >_{\text{lex}} xy^2 >_{\text{lex}} xy >_{\text{lex}} x >_{\text{lex}} y$$

Formal definition: $x^{\alpha} >_{\text{lex}} x^{\beta}$ if in the difference $\alpha - \beta$ (which belongs to $\mathbb{Z}^n$), the leftmost nonzero entry is positive.

$x^2yz^3 >_{\text{lex}} x^2z^4$ or $x^2z^4 >_{\text{lex}} x^2yz^3$?

$\rightarrow x^2yz^3 >_{\text{lex}} x^2z^4$ because $(2,1,3) - (2,0,4) = (0,1,-1)$
Monomial Orders on $k[x_1, \ldots, x_n]$ - the Graded Reverse Lexicographic Order $>_{\text{grevlex}}$

Let $x^\alpha$ and $x^\beta$ be monomials in $k[x_1, \ldots, x_n]$. $x^\alpha >_{\text{grevlex}} x^\beta$ if:

- $\sum_i^n \alpha_i > \sum_i^n \beta_i$, or if
- $\sum_i^n \alpha_i = \sum_i^n \beta_i$ and in the difference $\alpha - \beta$, the rightmost nonzero entry is negative.

Under this order $>_{\text{grevlex}}$:

$$xy^2 >_{\text{grevlex}} x^2 >_{\text{grevlex}} xy >_{\text{grevlex}} x >_{\text{grevlex}} y$$

$$x^2y^2z^2 >_{\text{grevlex}} xy^4z \quad \text{or} \quad xy^4z >_{\text{grevlex}} x^2y^2z^2 ?$$

$\rightarrow xy^4z >_{\text{grevlex}} x^2y^2z^2$ because $1 + 4 + 1 = 2 + 2 + 2$ and $(1, 4, 1) - (2, 2, 2) = (−1, 2, −1)$
Monomial Orders

\[ x^3 y^2 z >_{\text{lex}} x^2 y^6 z^8 \]
\[ x^2 y^6 z^8 >_{\text{grevlex}} x^3 y^2 z \]

\[ x^2 y^2 z^2 >_{\text{lex}} xy^4 z \]
\[ xy^4 z >_{\text{grevlex}} x^2 y^2 z^2 \]
Important property

If

▶ we use the monomial order $>_\text{lex}$ to compute a Gröbner basis and
▶ the solution set is finite,

then a univariate polynomial (in the last variable) is in the basis.

For example, the Gröbner basis for $\langle x^2 - y^2 + 1, xy - 1 \rangle$ is $\langle y^4 - y^2 - 1, x - y^3 + y \rangle$.

The system

$$\begin{align*}
x^2 - y^2 + 1 &= 0 \\
xy - 1 &= 0
\end{align*}$$

has the same solutions as the system:

$$\begin{align*}
y^4 - y^2 - 1 &= 0 \\
x - y^3 + y &= 0
\end{align*}$$

but the latter is much simpler to solve.
A More Ugly Example

A Gröbner basis for

\[
\begin{cases}
    x^2 - 2xz + 5 & = 0 \\
    xy^2 + yz + 1 & = 0 \\
    3y^2 - 8xz & = 0
\end{cases}
\]

under $>_\text{lex}$ is

\[
\{-81 + 4320z - 86400z^2 + 766272z^3 - 2513488z^4 - 295680z^5 - 242496z^6 + 61440z^8, -2472389942760 + 1450790919y + 98722479369600z - 1312504296363936z^2 + 5756399991700688z^3 + 711670127441280z^4 + 549519027506496z^5 - 10326680985600z^6 - 139421921341440z^7, 6503592729600 + 1450790919x - 257416379643438z + 3400639490020320z^2 - 14857079919551480z^3 - 1835782187164800z^4 - 1418473727285760z^5 + 26347944960000z^6 + 359882180198400z^7\}
\]